

Main Memory Operation Buffering for Efficient R-Tree Update Based on Biveinis, Šaltenis and Jensen (Aalborg University) paper in VLDB 2007

Setting

- A typical pervasive computing scenario
- Sampling of continuous processes via large numbers of sensors
- Maintenance of an up-to-date current state of the processes
- Query processing against the current state
- Example: moving objects
- The current positions of moving objects
- Large populations of objects are anticipated (mobile phone users)
- Updates are very frequent.
- E.g., 600,000 objects, each issuing an update every minute yields 10,000 updates per second.
- Indexing is essential to efficient query processing.
- The index must be stored on disk, at least in part.
- Existing indices do not support massive update loads.



Background: the R-tree

Observations

- Several updates to the same leaf cause separate traversals
- Update locality is not exploited
- Main memory is not used
- High rate of updates is required to sustain accuracy

The R^R-tree: the Data Structure

- A disk R-tree without any data structure modifications
- A single buffer for all incoming updates
- Any amount of memory, the more, the better
- Organized as a main-memory R-tree (can be accessed as a list too)
- Data contains the flag to tell deletions from insertions



The R^R-tree: Buffer Emptying

- When the buffer gets full, it is processed on the main tree •
- Updates travel down the tree, sharing I/Os •
- Small groups should be filtered





Experiments: Main-Memory Utilisation



- these to share I/O.
- existing proposals.
- Future work

- performance trade-off

Buffer size in objects

Summary

• Presented a new main-memory buffering technique for R-tree type indices

• The general idea is to speed up updates by allowing

• Uses partial buffer emptying

• Empirical studies show that the proposal improves on

• See the paper for the analytical study

• Application to other types of indices

• Better main-memory indexing

• Exploring the query performance/update